

Ditto: Efficient Serverless Analytics with Elastic Parallelism

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Serverless computing









Fine-grained resource elasticity



- Auto-scaling
- Concurrency from 1 to 1,000

Fine-grained billing



- 1 MB memory granularity
- 1 ms time granularity

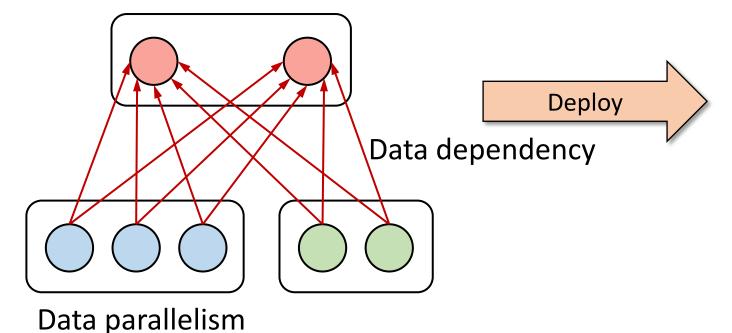
Serverless analytics



Big data & SQL-like query

Locus (NSDI'19) NIMBLE (NSDI'21) Databricks SQL Serverless Azure Synapse Analytics Google BigQuery

Job Execution DAG



Serverless functions

















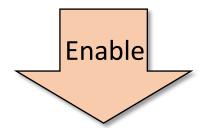
Job completion time (JCT)



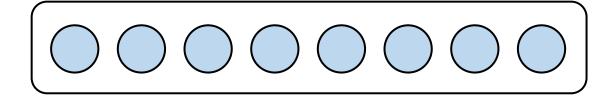
Cost (\sum_{funcs} time×memroy)

Degree of Parallelism: a new problem

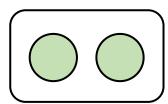
Fine-grained resource elasticity



Higher DoP Faster, lower JCT



Lower DoP
Lower cost

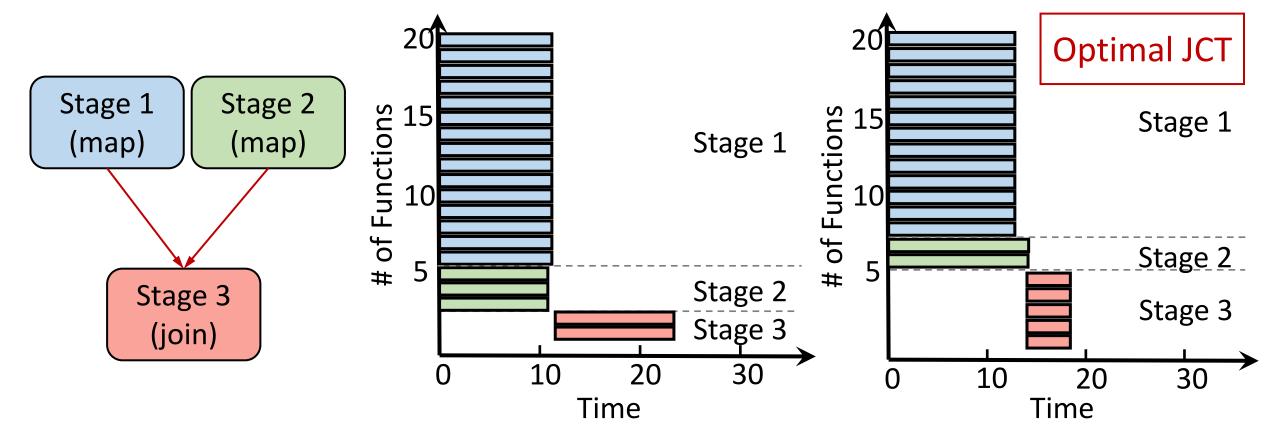


NIMBLE: a data perspective

DoP proportional to input data size

Caerus: NIMBLE Task Scheduling for Serverless Analytics





Elastic parallelism



Data? X
Time? ✓







Job completion time (JCT)



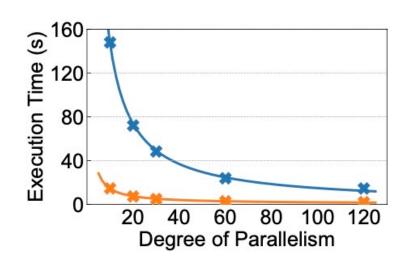
Cost ($\sum_{\text{funcs}} \text{time} \times \text{memroy}$)

Main idea:

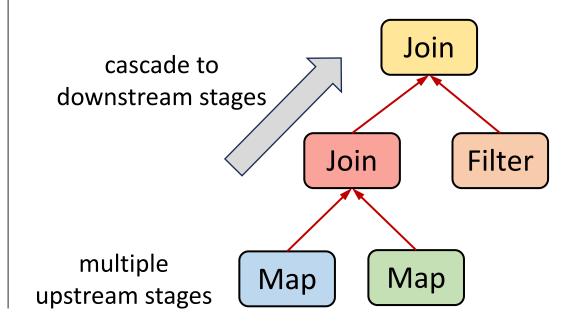
- Match the resource elasticity of serverless computing with parallelism scheduling in data analytics
- Optimize serverless performance goals directly from a perspective of time

Challenge 1: Optimal parallelism for arbitrary DAGs

 Accurate prediction of the execution time under dynamic parallelism configurations



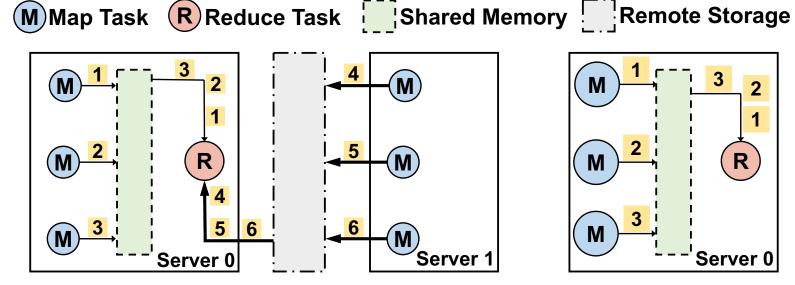
Consider data dependencies

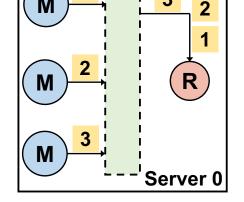


Challenge 2: Coupling of parallelism and placement

Co-optimize parallelism configuration and function placement

Shared memory SPRIGHT (SIGCOMM'22) Pheromone (NSDI'23)





High DoP with heavy data shuffle time

Low DoP with almost zero data shuffle time

Ditto design outline

Challenge 1: How to find the optimal parallelism for arbitrary DAGs?

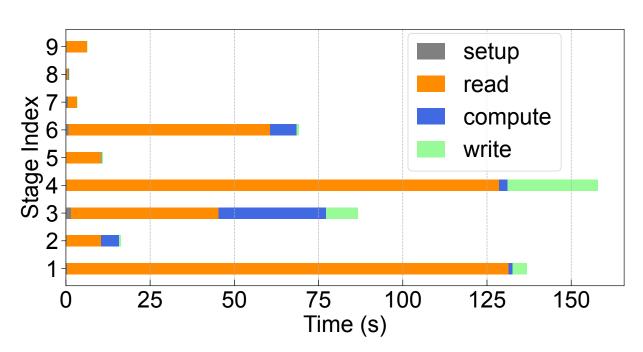
- Execution time model → Time under dynamic parallelism
- DoP ratio computing → Optimal parallelism configuration

Challenge 2: How to optimize the coupled parallelism and placement?

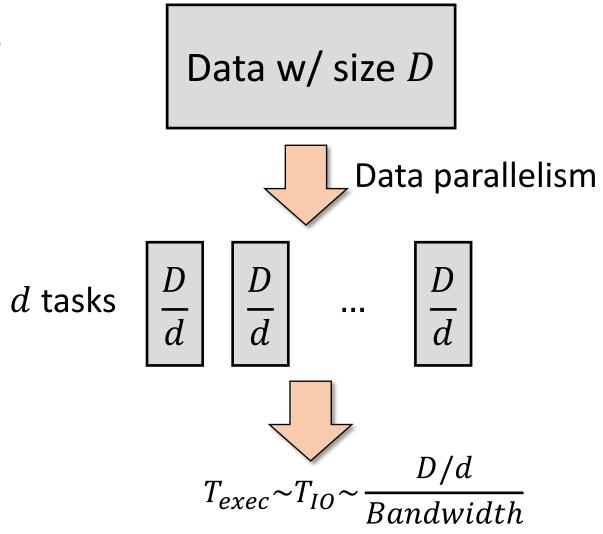
- Greedy grouping → Eliminate high data shuffling overhead
- Joint iterative optimization → Co-scheduling

Execution time model: a time perspective

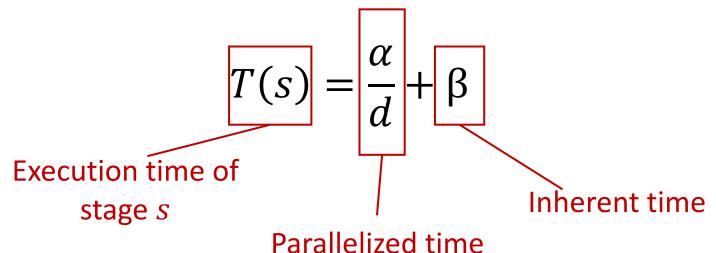
- Long running: 10 to 1000 seconds
- Data I/O dominates



Time breakdown for TPC-DS Q95



Execution time model: a time perspective



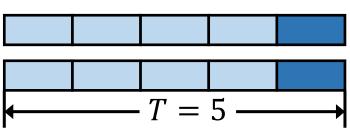
d: degree of parallelism, DoP

 α : the parallelized time parameter

$$\alpha = 8$$
, $\beta = 1$

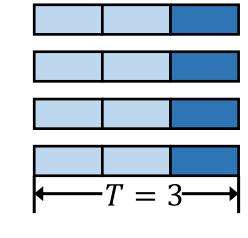


Inherent time unit



$$T=5$$

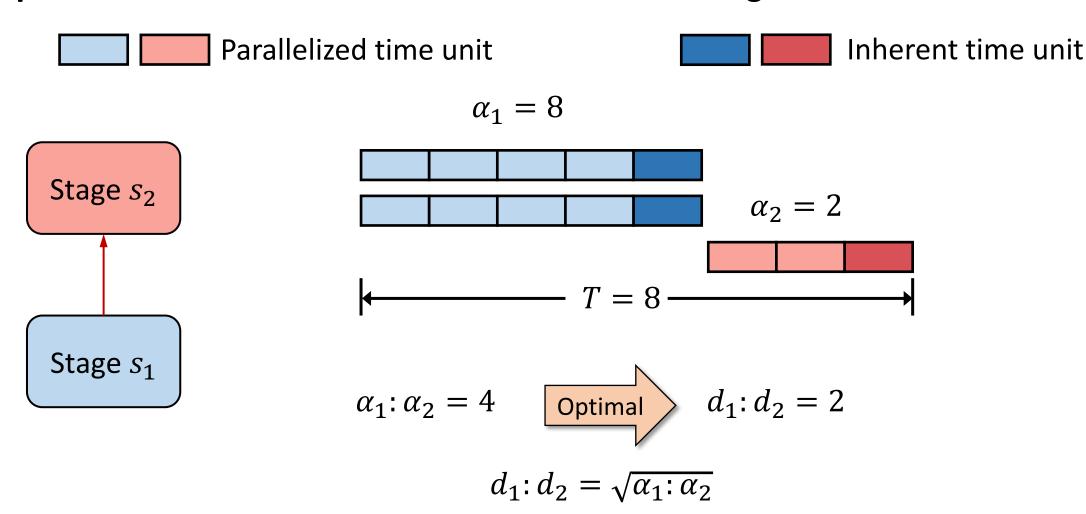
$$d=2$$



$$d = 4$$

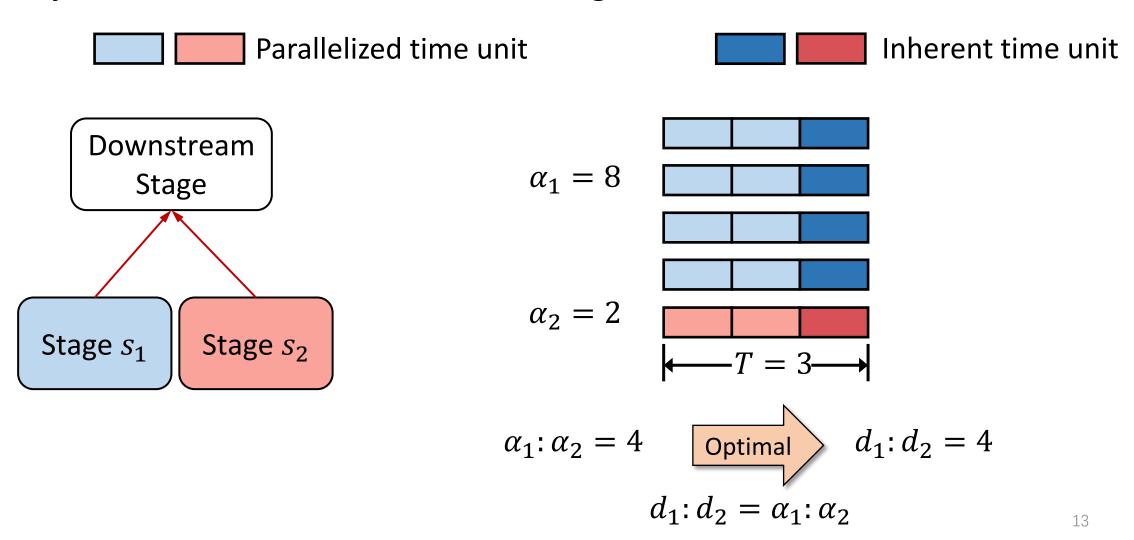
DoP ratio computing

Intra-path DoP ratio: minimize the sum of the two stages' execution time



DoP ratio computing

Inter-path DoP ratio: balance the two stages' time

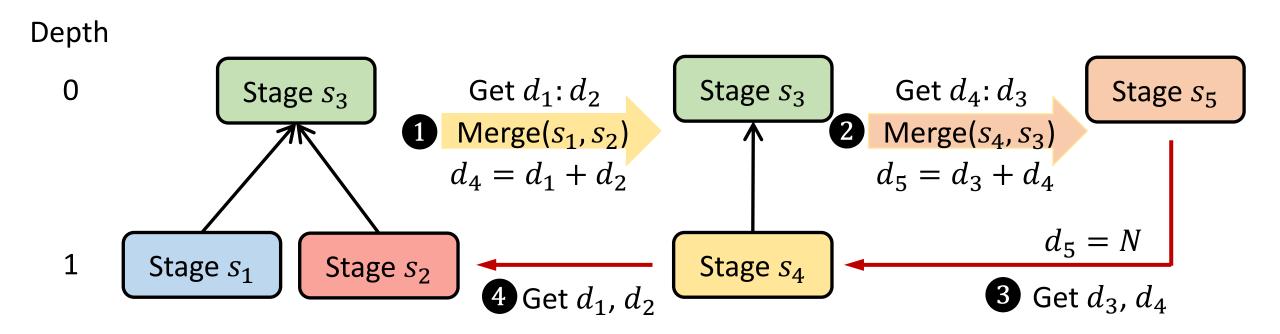


DoP ratio computing

Stage merging: a new stage also conforms to the execution time model

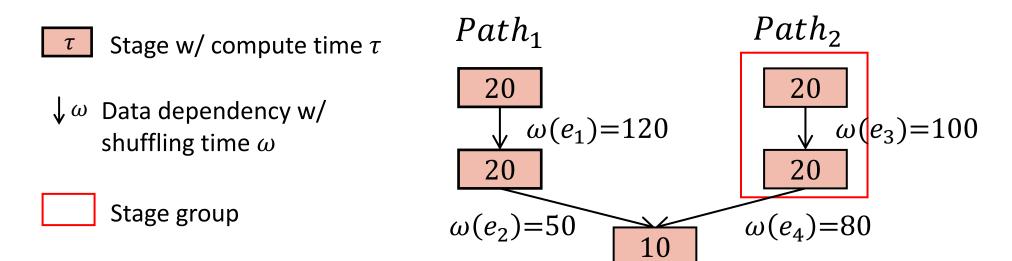
 d_i : degree of parallelism of stage s_i

N : total number of functions



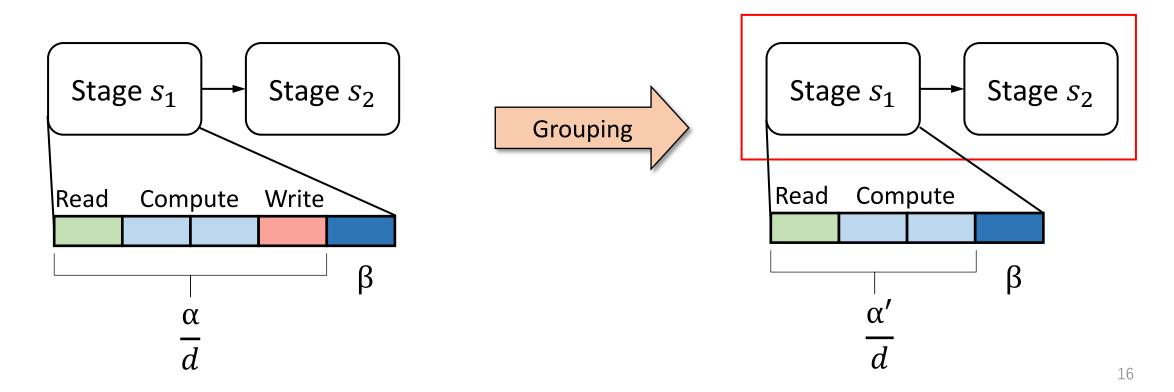
Greedy grouping

- Stage group: stages that should communicate via shared memory
 - NP-hard
- Greedy order: group stages with high shuffling overhead
 - For JCT optimization, the highest on the critical path first

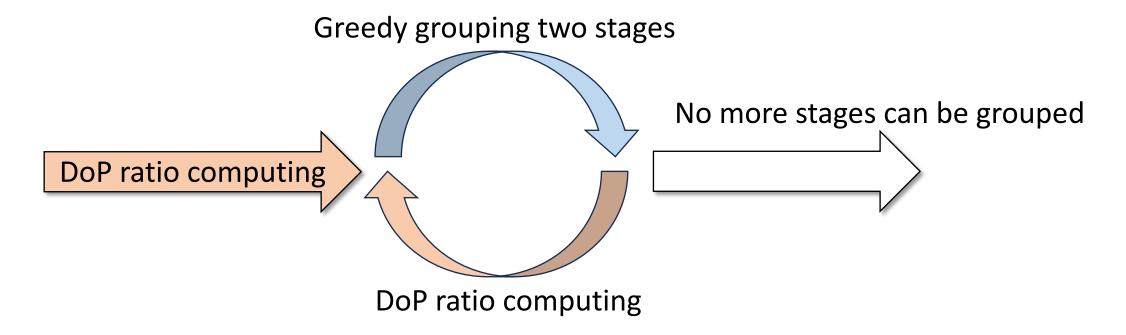


Joint iterative optimization

- α will decrease as the I/O time reduces to zero after grouping
 - Model the I/O and compute parts of α separately
 - Combine with DoP ratio computing into joint optimization



Joint iterative optimization



- Each stage is a group initially
- In each iteration
 - group two stages (or stage groups) with the highest shuffling overhead
 - recalculate the new optimal parallelism configuration

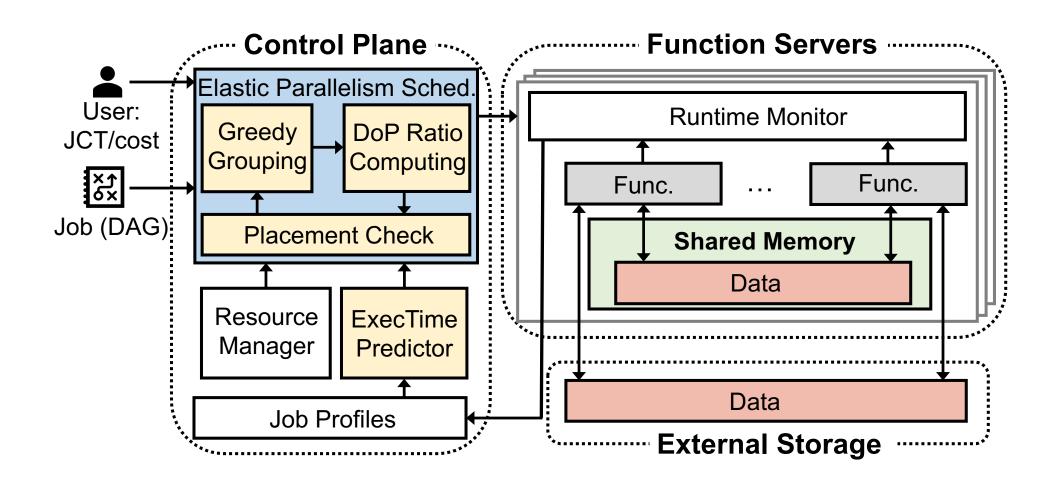
Cost optimization

- DoP ratio computing applies serverless cost model
 - Function cost: consider the resource usage
 - Total cost: the sum of all function costs
- Greedy grouping groups stages with highest shuffling cost first

Please refer to our paper for more details!

Ditto System

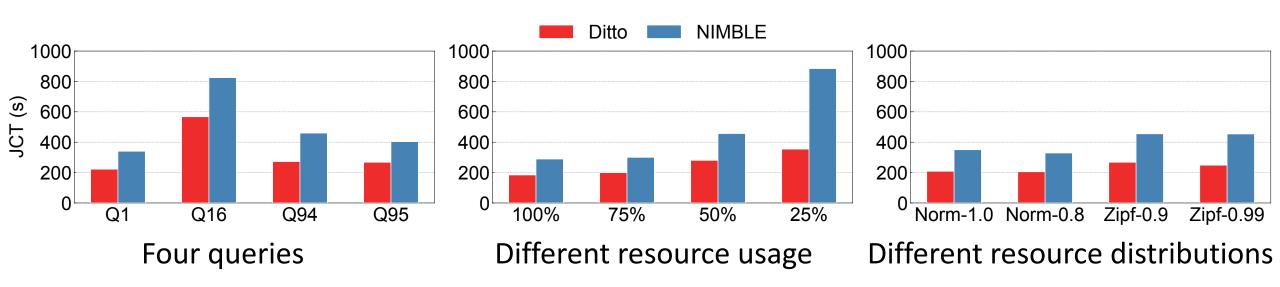
Implement Ditto on top of SPRIGHT (SIGCOMM' 22)



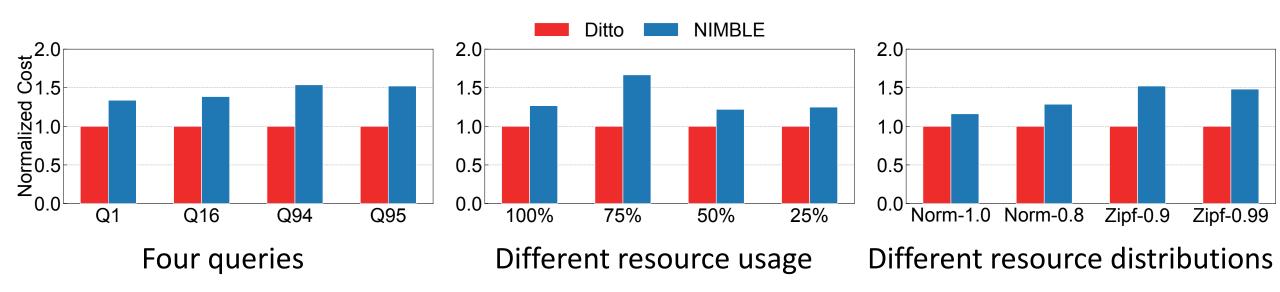
- Setup on AWS
 - Scheduling: one m6i.4xlarge server
 - Compute: eight m6i.24xlarge servers (96 vCPUs & 384 GB DRAM each)
 - Storage: S3
- TPC-DS
 - Q1, Q16, Q94, Q95
 - groupby, filter, join
 - 1 TB data

```
select
  count(distinct ws_order_number) as "order count",
  sum(ws_ext_ship_cost) as "total shipping cost",
  sum(ws_net_profit) as "total net profit"
from
 web_sales ws1,
 date_dim,
  customer_address,
 web_site
where
  d_date between '1999-4-01'
  and (cast('1999-4-01' as date) + 60 days)
  and ws1.ws_ship_date_sk = d_date_sk
  and ws1.ws_ship_addr_sk = ca_address_sk
  and ca state = 'IA'
  and ws1.ws_web_site_sk = web_site_sk
  and web company name = 'pri'
```

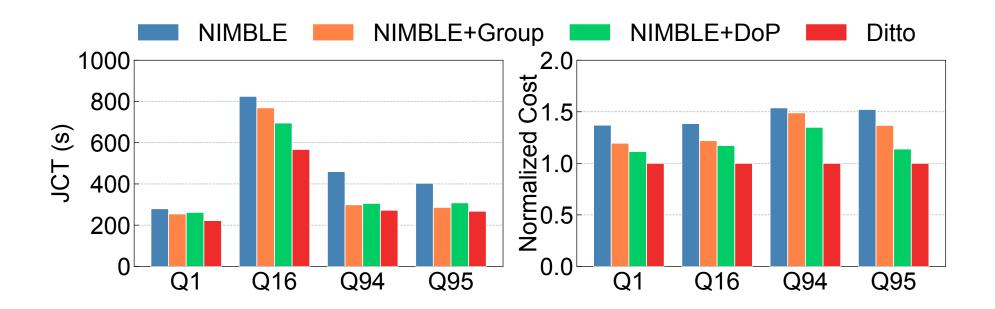
Ditto reduces the JCT by 1.3-2.5X compared to NIMBLE



Ditto reduces the cost by 1.2-1.7X compared to NIMBLE



Ablation experiment to verify the effectiveness of Ditto



- Performance under Redis
- Accuracy of the execution time model
- Execution breakdown for TPC-DS Query 95
- System overhead of Ditto

Conclusion

- Serverless analytics introduces the elastic parallelism scheduling problem to optimize serverless performance goals, i.e., JCT and cost
- Ditto co-optimizes parallelism configuration and function placement from the perspective of time
 - Execution time model under dynamic parallelism
 - DoP ratio computing to achieve optimal JCT or cost
 - Joint iterative optimization for both parallelism and placement
- Ditto reduces up to 2.5X in JCT and up to 1.7X on cost compared to NIMBLE

